

Carrier-Grade NAT

Josef Ungerman

Consulting Systems Engineer, Cisco, CCIE#6167

Email, Jabber, TP: josef.ungerman@cisco.com

Agenda

CGN (Carrier-Grade NAT)

Definition and purpose

NAT vs. Firewall

Mapping/filtering, app traversal

Other CGN Behaviors

Pooling, port limits, etc.

Session Logging

Netflow/Syslog, formats, variations

CGN Design

Performance, placement

Network Address & Port Translation

Most of Broadband users are behind NAT today!

When say “NAT”, they typically mean “NAPT”

- NAT

First described in 1991 (draft-tsuchiya-addrtrans), RFC1631

1:1 translation: Does not conserve IPv4 addresses

Per-flow stateless

Today's primary use is inside of enterprise networks

Connect overlapping RFC1918 address space

Note: NAT66 is stateful or stateless, but it is not NAPT

- NAPT

Described in 2001 (RFC3022)

1:N translation

Conserves IPv4 addresses

Allows multiple hosts to share one IPv4 address

Only TCP, UDP, and ICMP

Connection has to be initiated from 'inside'

Per-flow stateful

Commonly used in home gateways and enterprise NAT

“NAT44” is used to differentiate IPv4-IPv4 NAPT from Address Family Translation, typically referred to as NAT64 and NAT46”

What is CGN?

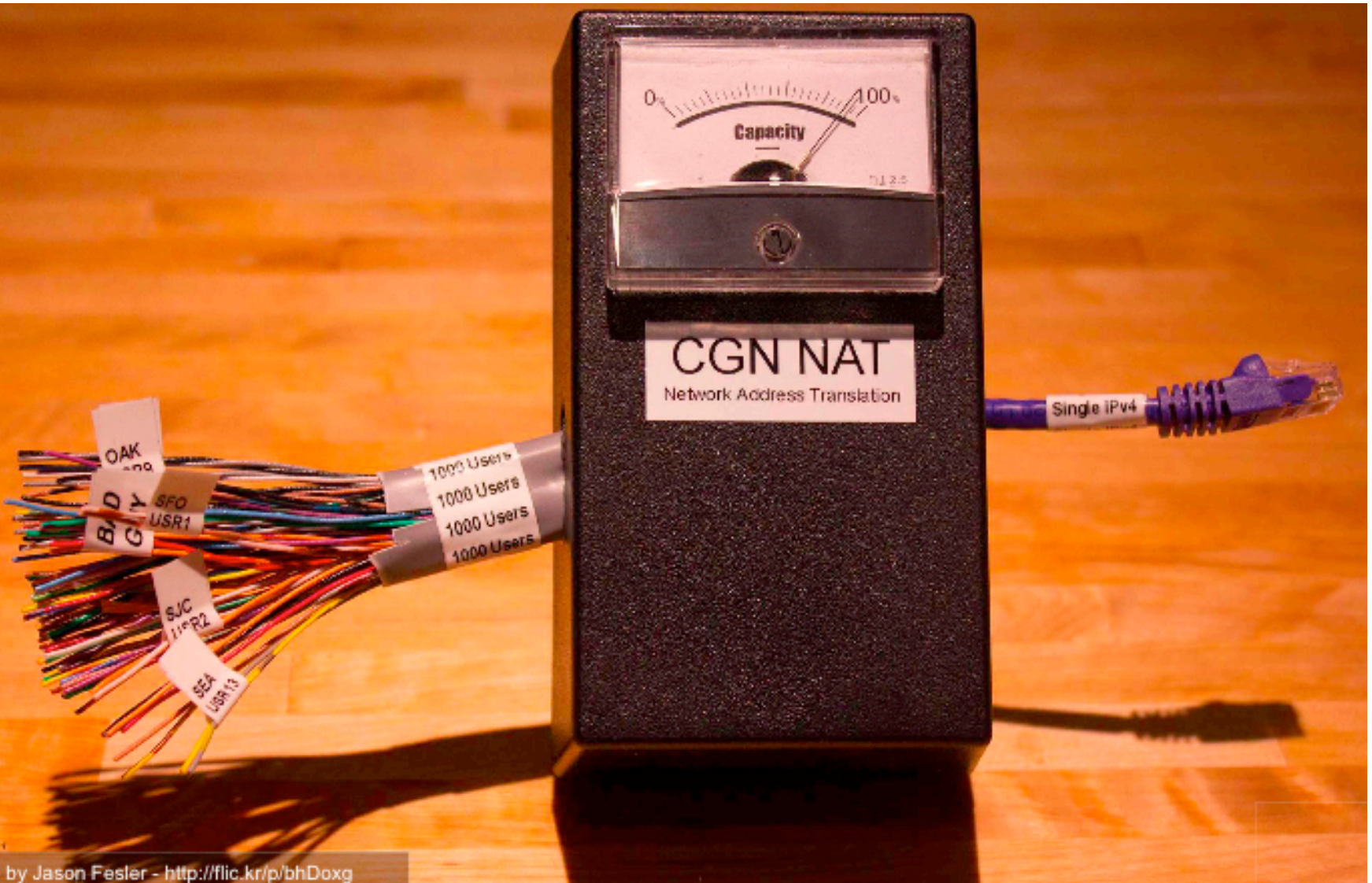
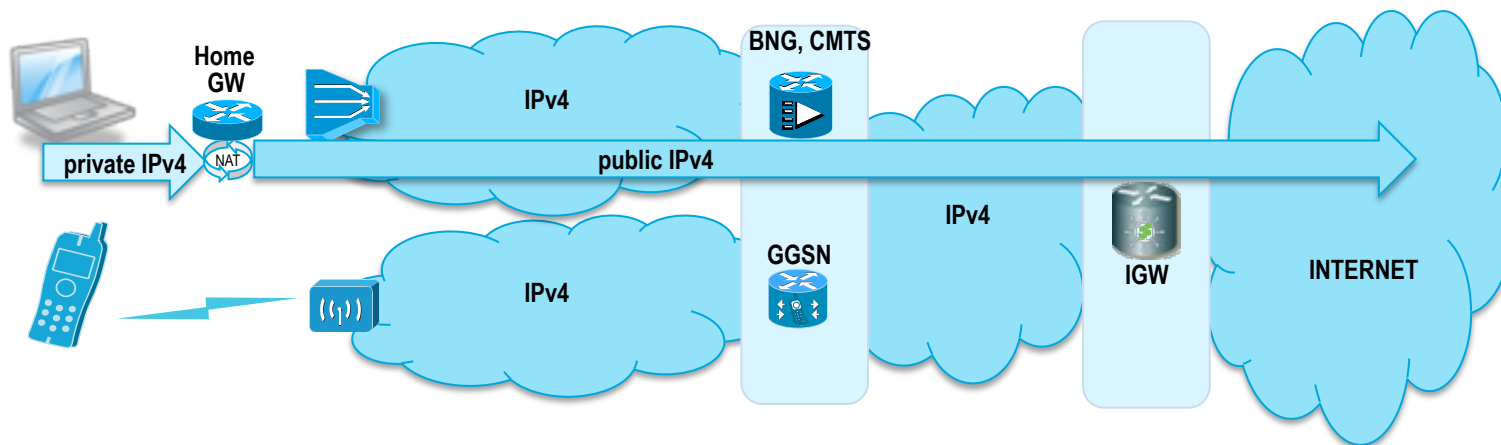


Photo by Jason Fesler - <http://flic.kr/p/bhDoxg>

Courtesy of Jason Fesler, Yahoo (V6 World Congress 2012)

NAT in Internet Access

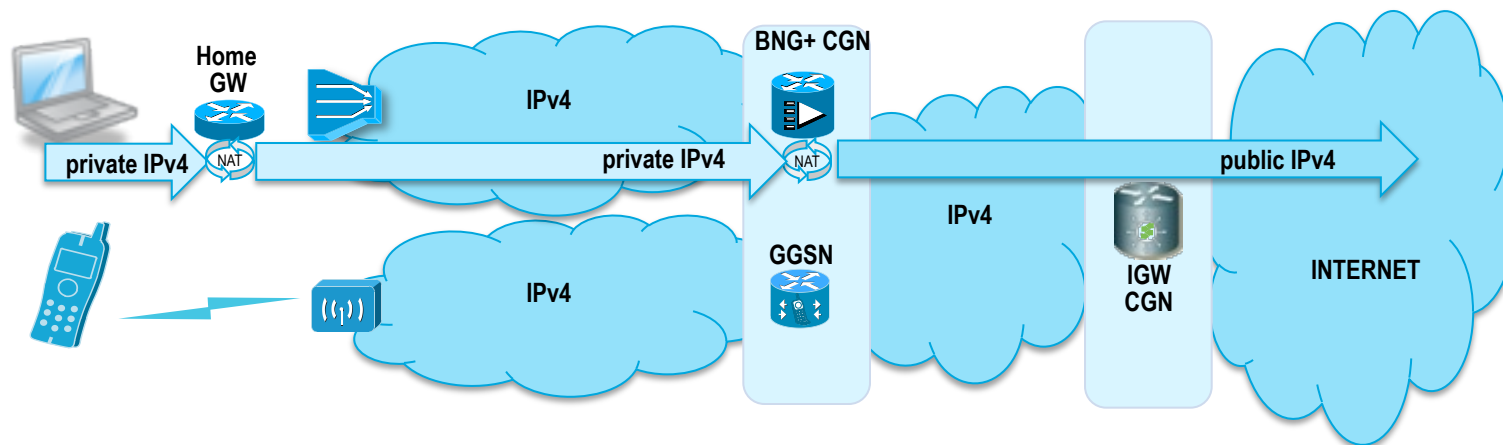
typical deployment today (wireline)



NAT in Internet Access

CGN – NAT444 (wireline)

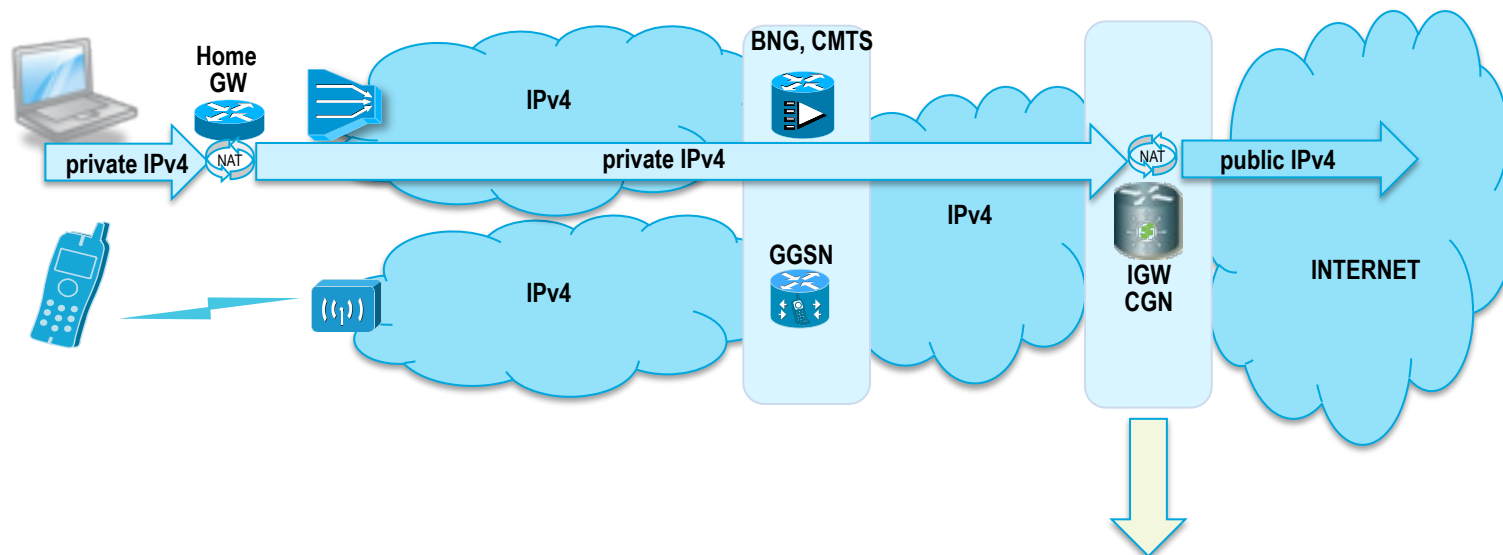
RFC1918 or RFC6598 (100.64.0.0/10)



NAT in Internet Access

CGN – NAT444 (wireline)

RFC1918 or RFC6598 (100.64.0.0/10)

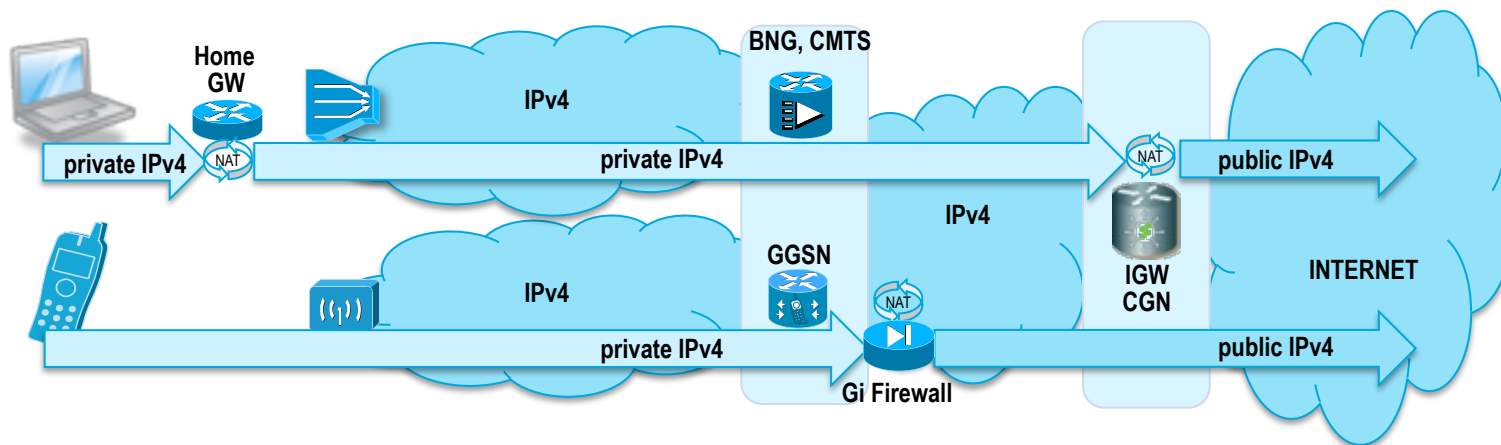


Large Scale

- 100G+ throughput
- 100M+ concurrent bidirectional sessions
- 1M+ sessions per second setup rate

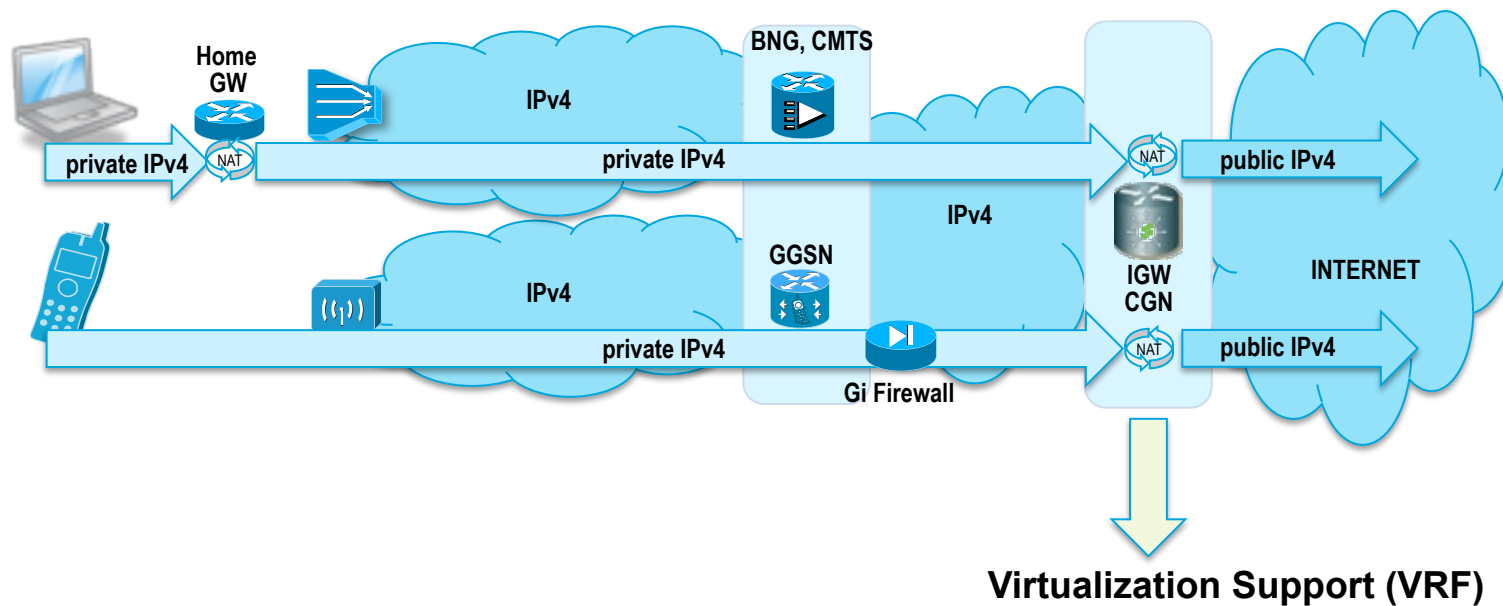
NAT in Internet Access

Mobile



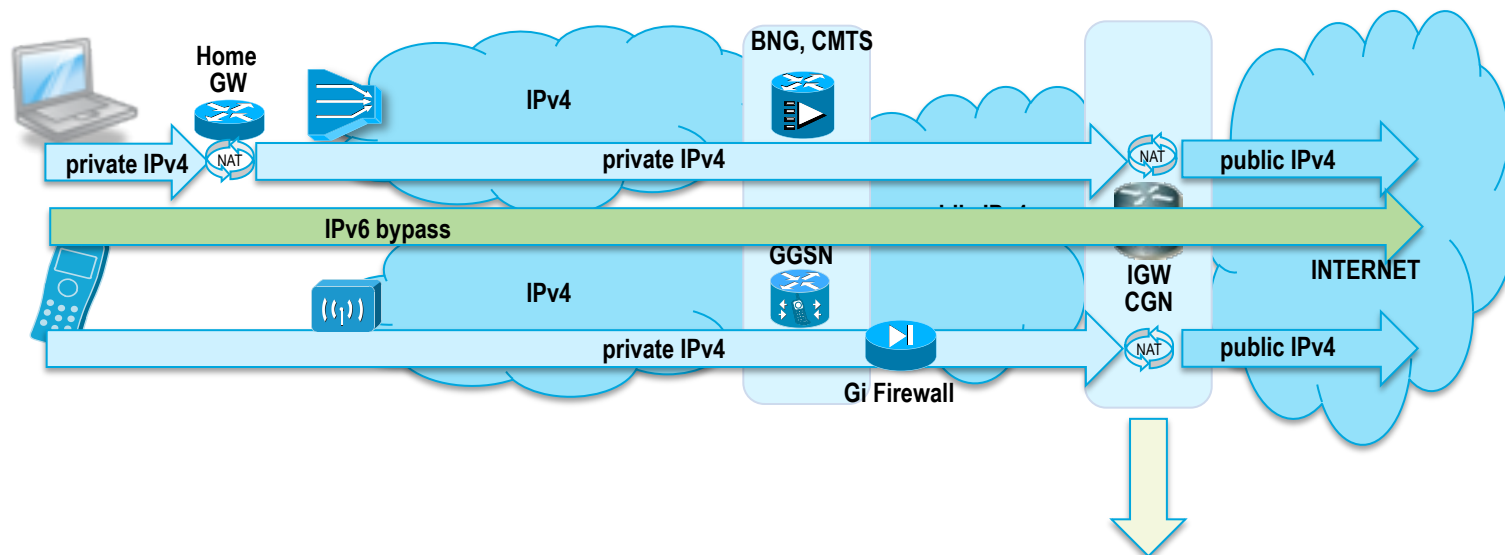
NAT in Internet Access

Mobile



NAT in Internet Access

IPv6 ultimately bypasses CGN



Dual-Stack
6in4 tunneling – 6rd BR
4in6 tunneling – DS-Lite AFTR, MAP
NAT64 for v6-only hosts

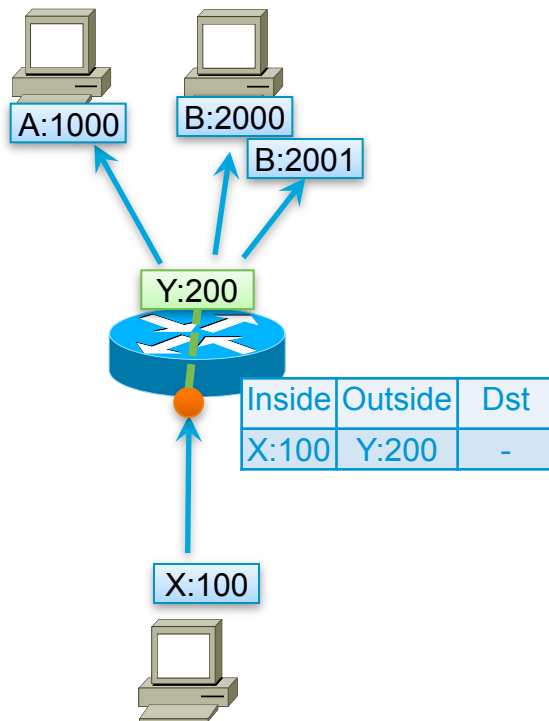
CGN = IP Address Sharing

IPv4 Exhaust workaround: move part of Internet from L3 to L4

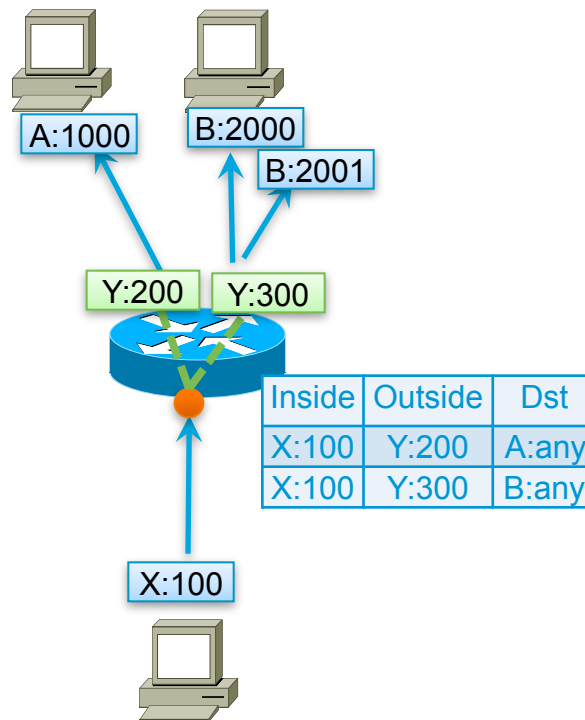
- Adds capital and operations cost [\$/Gbps] → decreases ARPU
CGN is here not to stay, but to be replaced (by IPv6)
- Inherent issues (general NAPT issues, mostly not NAT444 related)
[draft-ford-shared-addressing-issues](#)
- Servers must log also source port numbers
Shared IP address = shared suffering (blacklisting, spam,...)
Tracking and Law Enforcement, [draft-ietf-intarea-server-logging-recommendations](#)
Otherwise CGN must do log also Destination IP:port (privacy issue!)
- Requesting specific ports – “Not everyone can get port 80”
- Geo-Location issues (“get me the nearest ATM”)
- Complicates inbound access to media
- Keepalives → power consumption, mobile battery drain

NAPT internals: Mapping

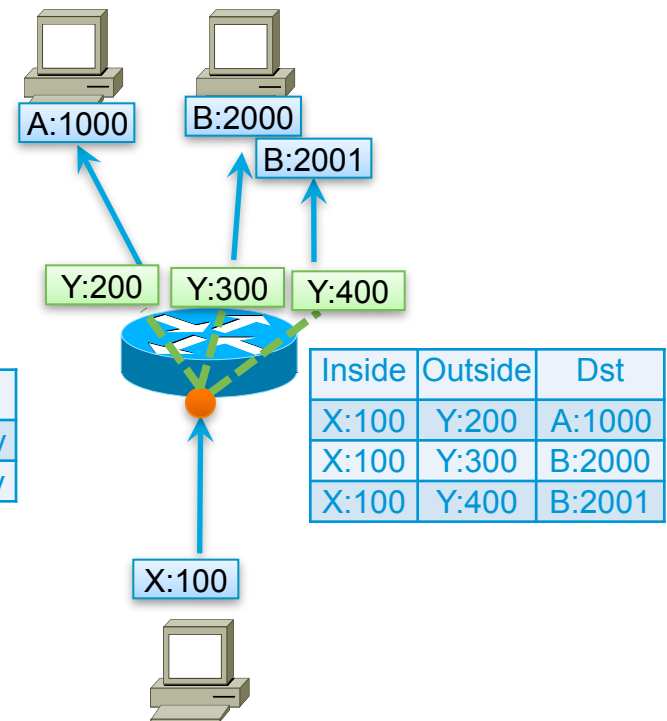
Endpoint Independent



Address Dependent



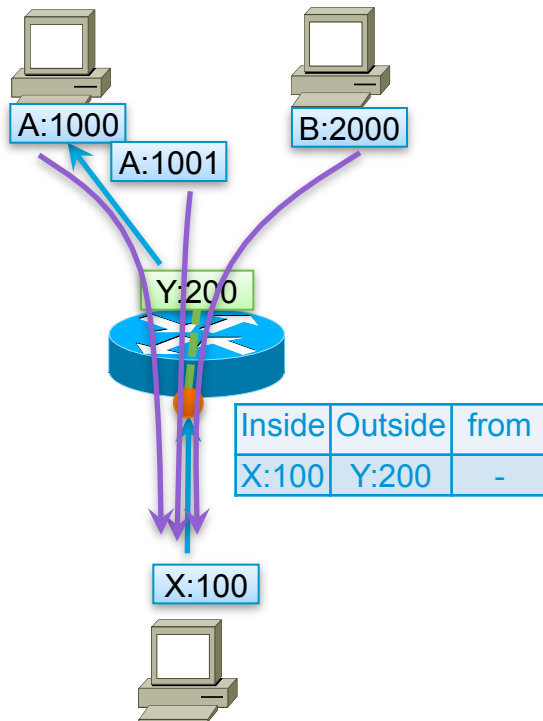
Address and port Dependent



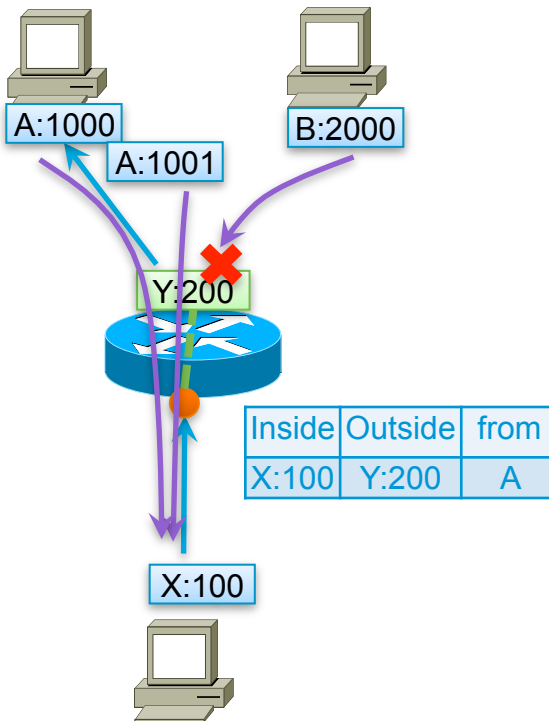
IP Address: Port Number

NAPT internals: Filtering

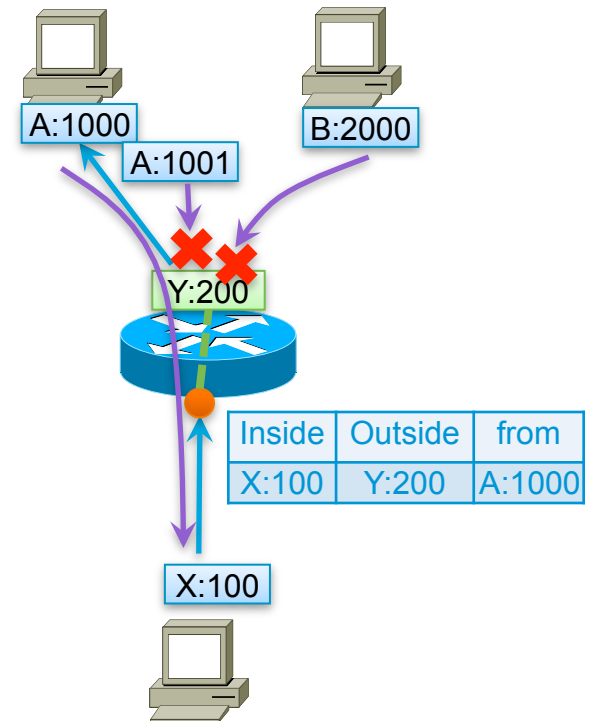
Endpoint Independent



Address Dependent

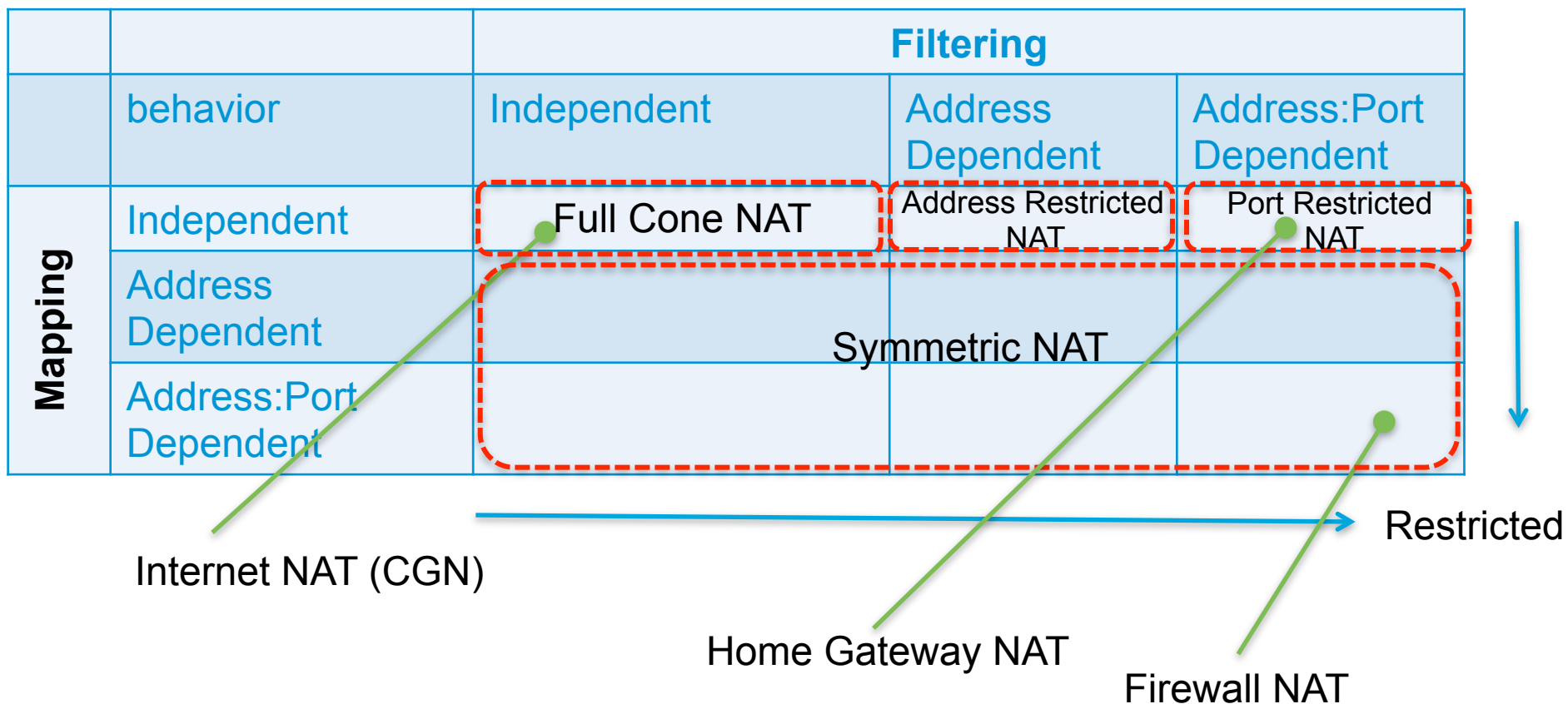


Address and Port Dependent



IP Address: Port Number

NAT mapping/filtering behavior

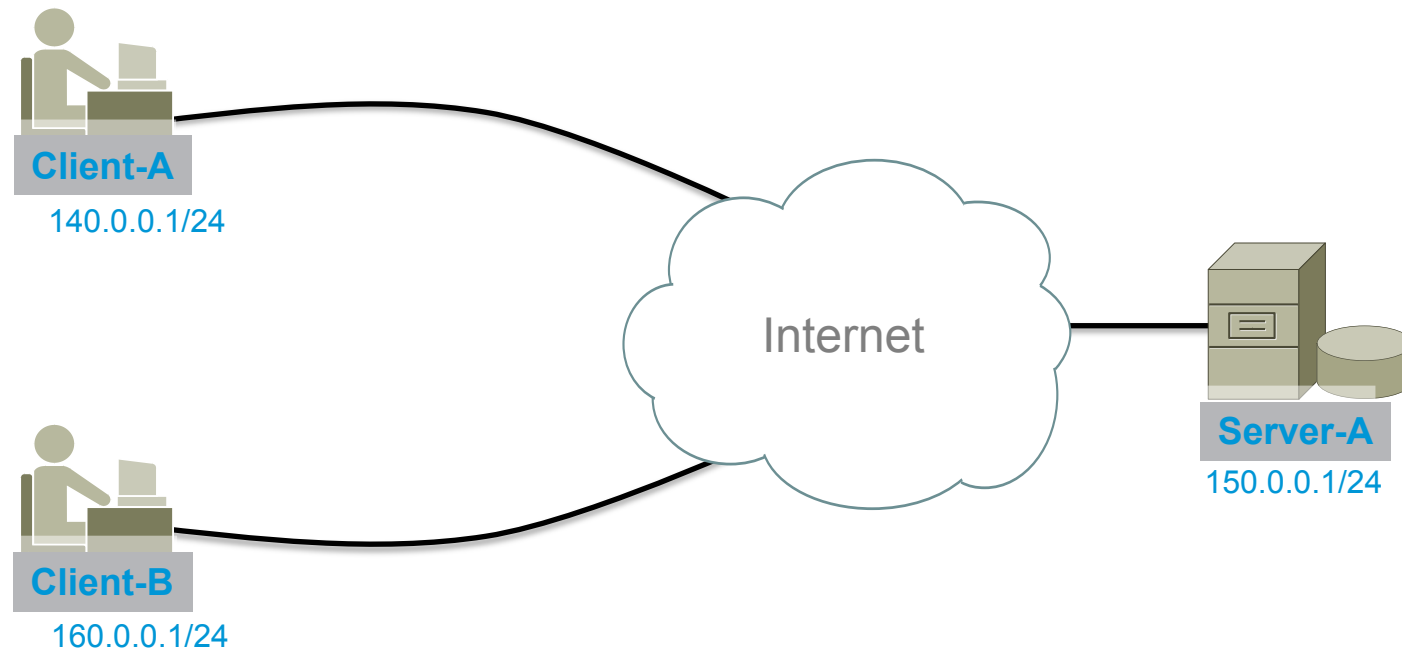


STUN NAT Types

- Classic STUN : simple traversal of UDP through NAT(RFC3489)
- Now : Session Traversal Utilities for NAT(RFC5389)

NAT != Firewall

Application transparency behaviors

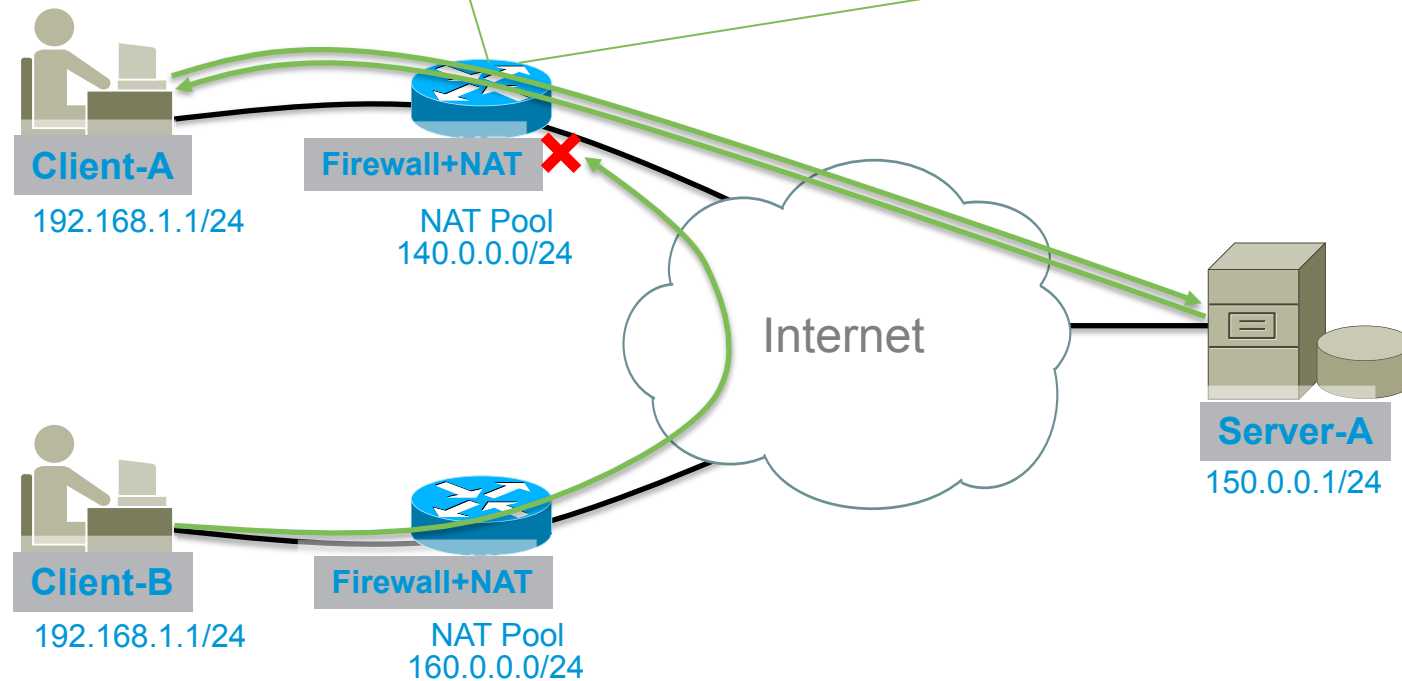


What is Symmetric NAT

Address:port restricted NAT = Firewalling behavior

NAT/PAT Entry:

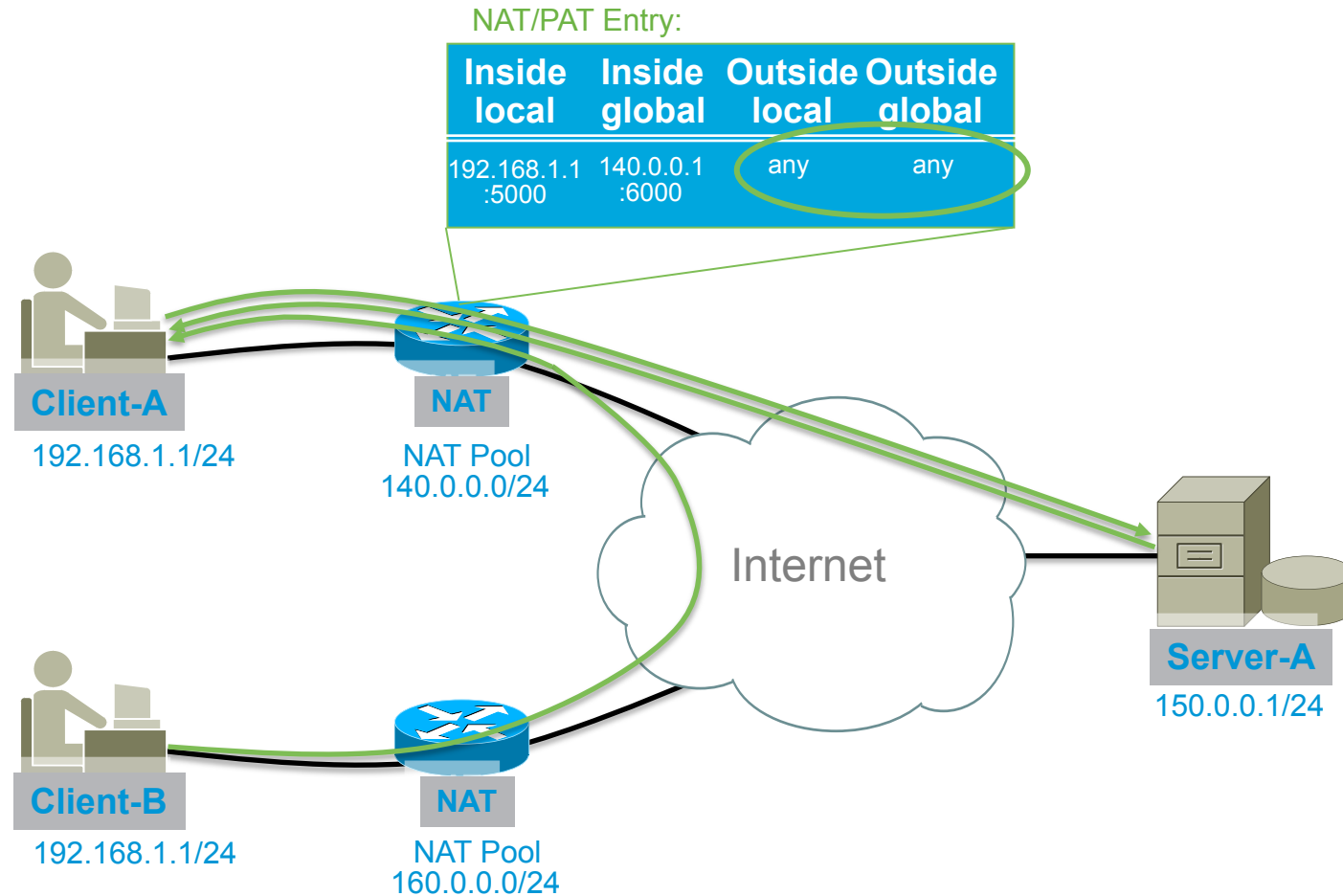
Inside local	Inside global	Outside local	Outside global
192.168.1.1 :5000	140.0.0.1 :6000	150.0.0.1 :6000	150.0.0.1 :6000



Endpoint Dependent Mapping, Endpoint Dependent Filtering

What is Full cone NAT

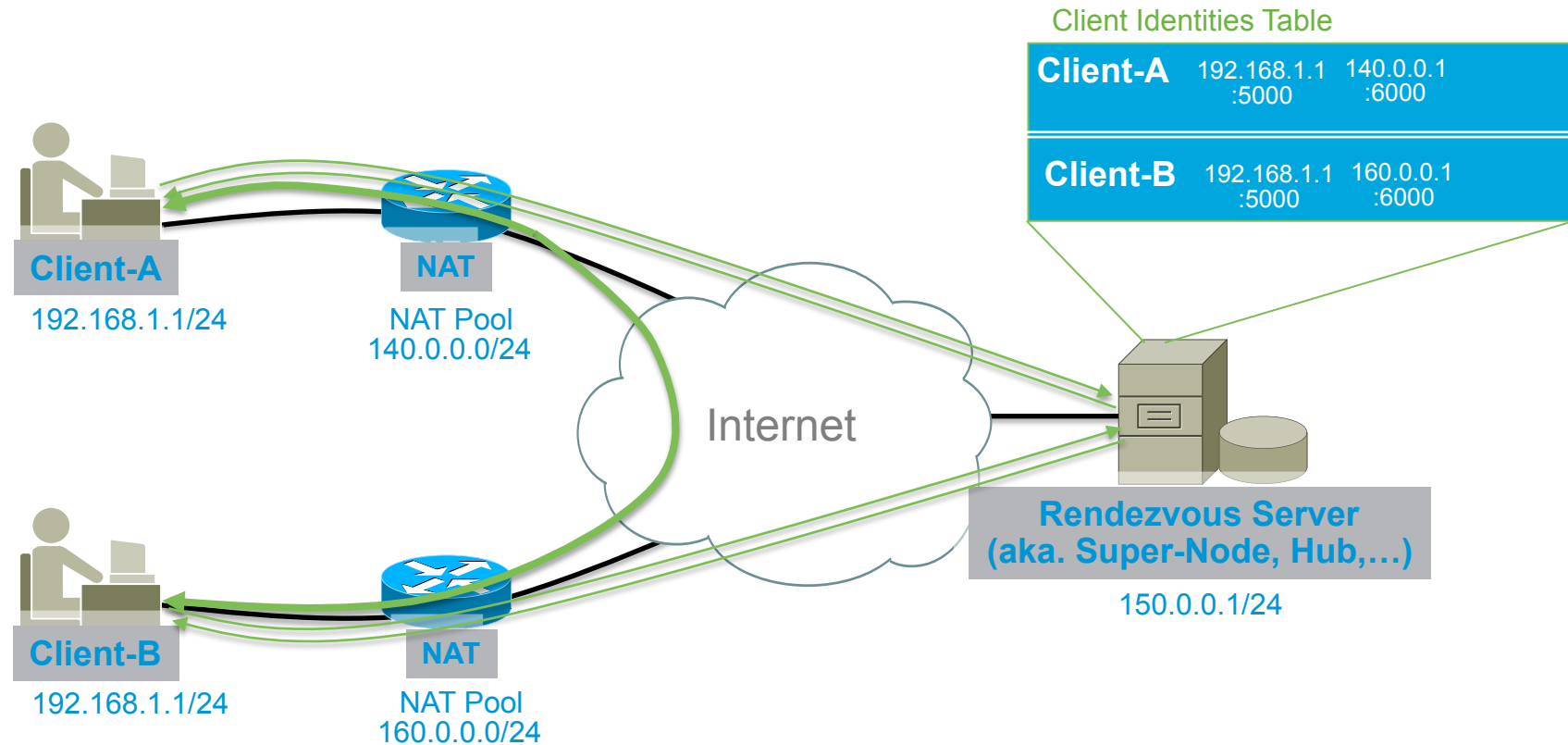
Pure NAT with no Firewalling behavior



Endpoint Independent Mapping, Endpoint Independent Filtering – EIM/EIF

NAT Traversal with Full cone NAT

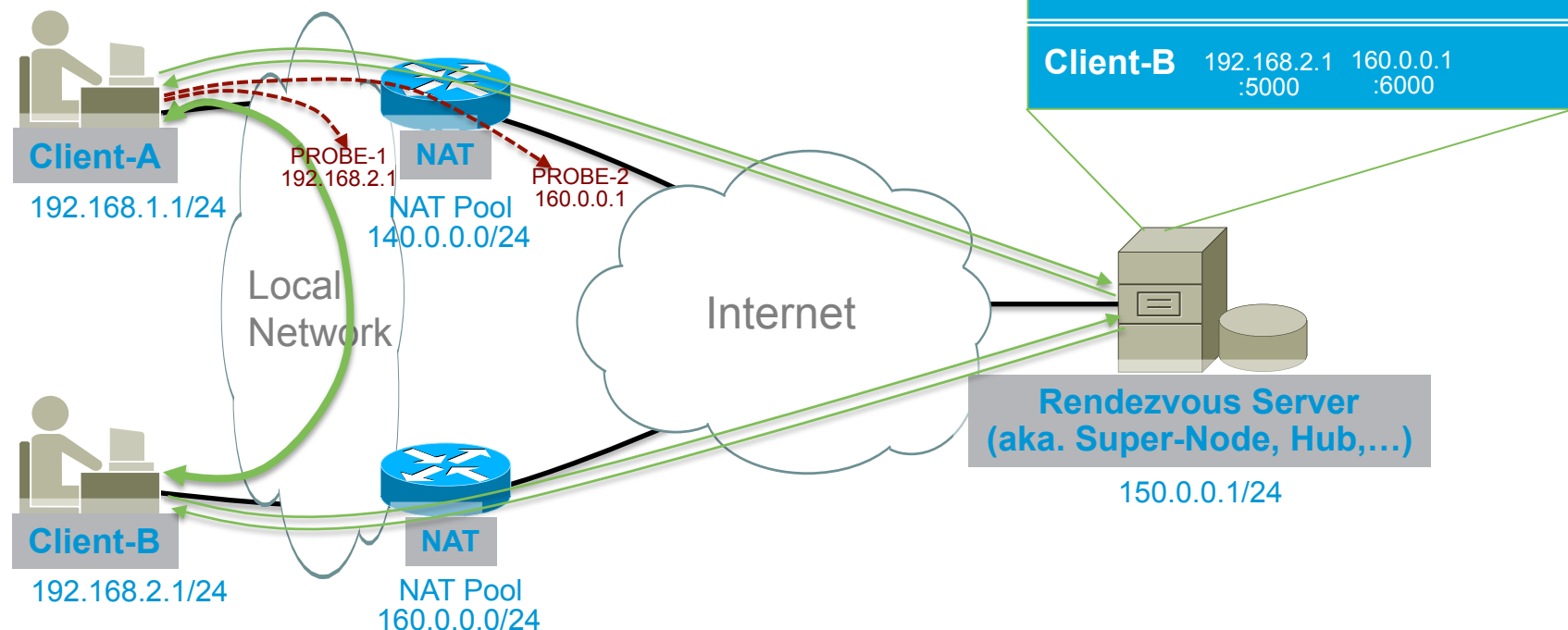
Peer-to-peer Applications Transparency



NAT Traversal with Full cone NAT

Connectivity Probes and the Shortest Path

One of the PROBE replies comes back first
(probably on-net one → fastest path)

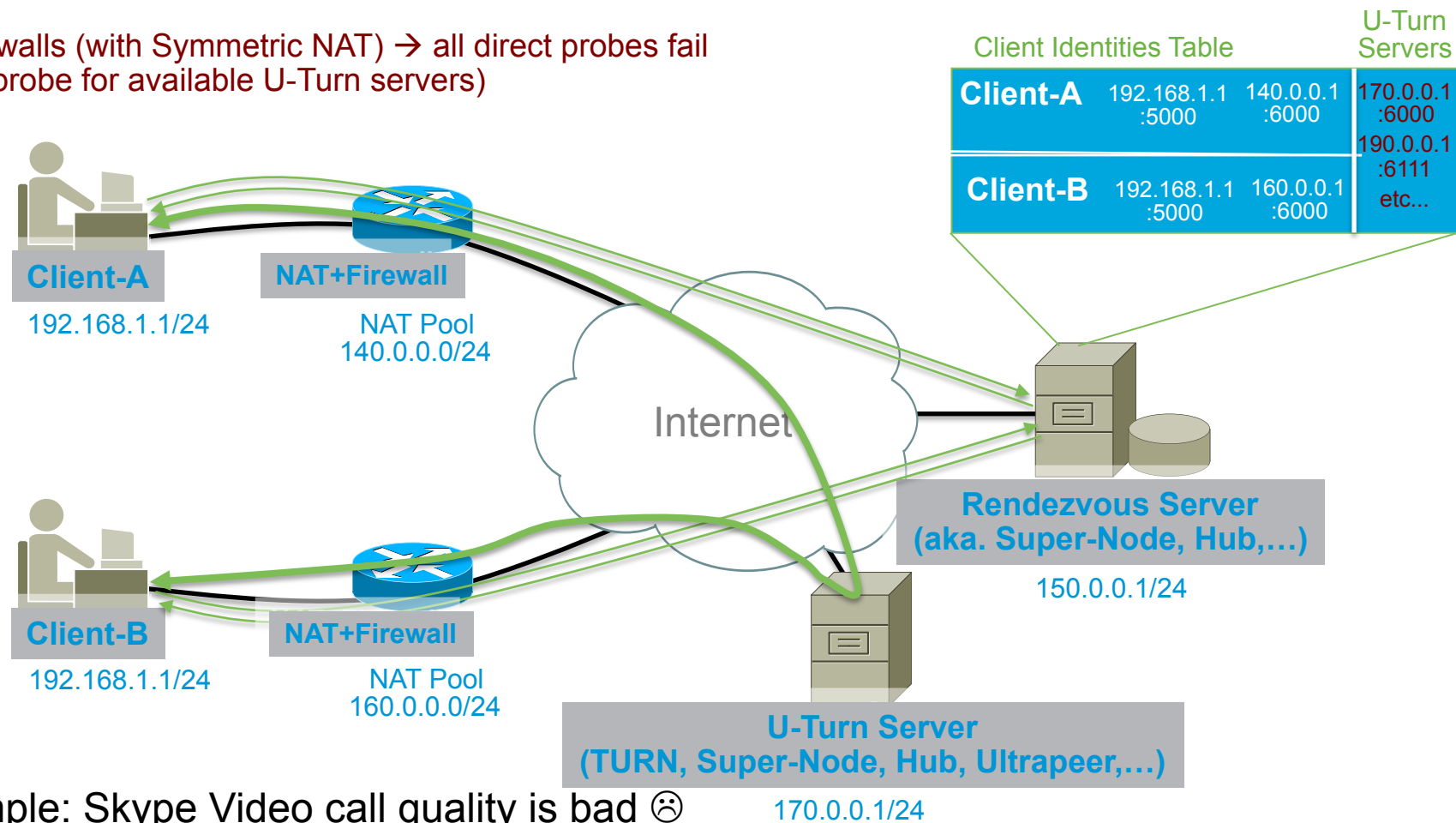


Example: Skype Video call quality is good 😊

NAT Traversal with Firewalls

U-Turn server (Super-Node)

Firewalls (with Symmetric NAT) → all direct probes fail
(→ probe for available U-Turn servers)



Example: Skype Video call quality is bad ☹

170.0.0.1/24

Avoiding U-Turn Servers with Firewalls

Opening port holes in Firewall/NAT

- Application talks to the Firewall

 - UPnP-IGD (app to home router)

 - Common in Home Environment

 - Not applicable in Internet (scale, no suitable protocols defined)

- ALG's (Application Level Gateway) aka. Fixup

 - App proxy in the firewall – we jump from L4 to L6+

 - Common in Enterprise (closed environment – we know the list of supported apps)

 - Hardly applicable in Internet (open environment – we don't know all the apps)

 - How about encrypted and integrity-protected protocols?

 - How about URL with literals <https://1.2.3.4>

Summary: CGN behavior best practices

draft-behave-lsn-requirements, RFC4787, RFC6145

- Internet way: keep it simple
 - EIM/EIF → Full Cone NAT → no ALG's (with well controlled exceptions)
 - CGN role is IPv4 exhaust solution, not security, LI, traffic monitoring, etc.
 - Respect OSI model – stay at L3, at L4 if you have to, not higher layers
- Firewall = ALG's = going above L4
 - Breaks Net Neutrality (OTT regulatory pushback)
 - ALG for Vendor-A breaks app from Vendor-B (same port, different traversal)
 - Undefined performance impact of ALG's → numerous DoS attack vectors
 - Bugs, ISP liable for 3rd party apps?

Reality Check: Internet Apps work with NAT



iTunes



Google Maps



Playstation Network



Google Chrome



iPhone App Store



Windows Live Messenger



Google Talk



Reality Check: Apps and NAT Traversal

- STUN, ICE, TURN

NAT EIM/EIF – Intelligence in endpoint

Useful for offer/answer protocols
(SIP, XMPP, probably more)

Standardized in MMUSIC and BEHAVE

ICE apps – examples:

- Google chat (XMPP)
- Microsoft MSN (SIP inside of XML)
- Yahoo (SIP)
- Counterpath softphone (SIP)

STUN: “Session Traversal Utilities for NAT” – RFC 5389

ICE: “Interactive Connectivity Establishment” – RFC 5245

TURN: “Traversal Using Relays around NAT” – RFC 5766

- Other examples

IPSec over TCP/UDP

FTP PASV – data connection always to server

RTSPv1 → RSTPv2 (effectively replaced HTTP Video, ABR,...)

Skype – encrypted, does its own NAT traversal

Port 80/443 apps

Known Problems:

Active FTP (old browsers), RTSPv1 (old m.youtube.com), MS PPTP (old PC VPN)

CGN Behaviors

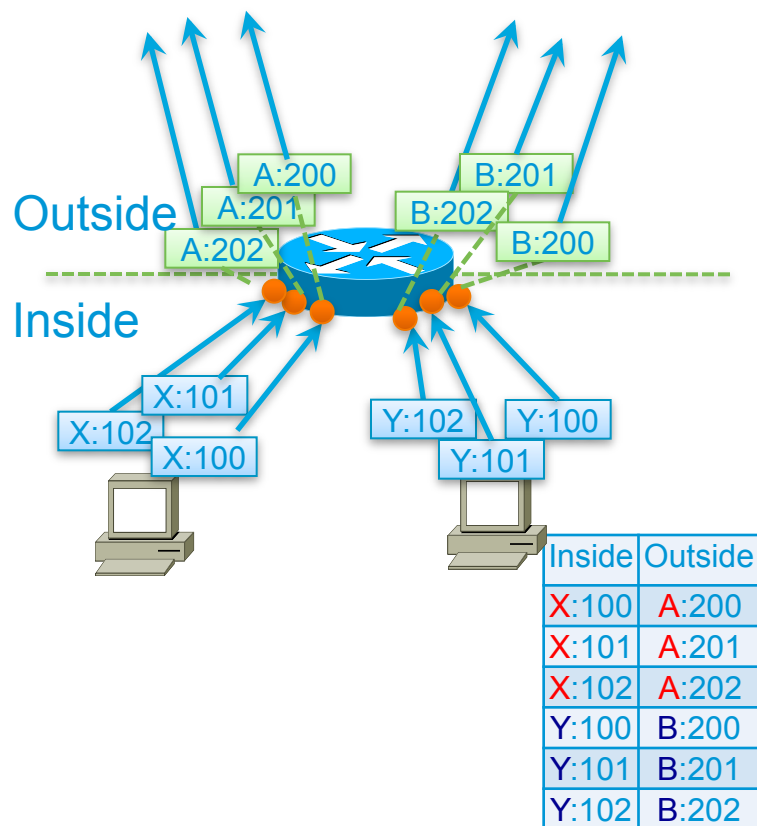
draft-ietf-behave-lsn-requirements



A CGN is defined by constrained behavior:

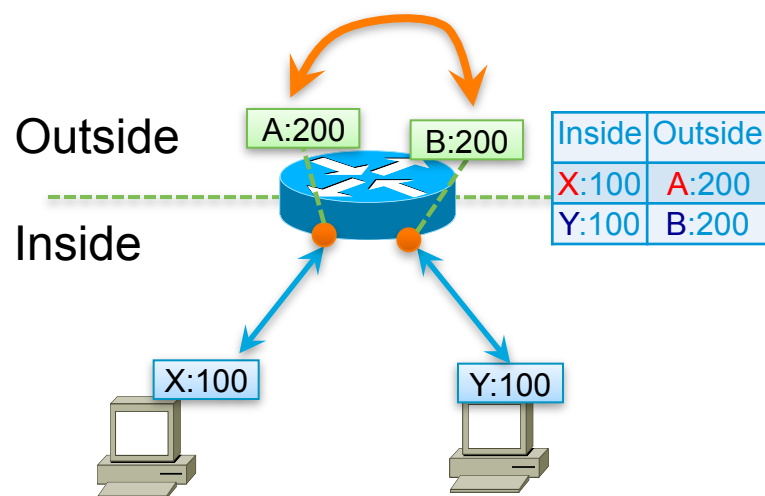
- NAT Behavior Compliance ([RFC4787](#), [RFC5382](#), [RFC5508](#))
 - Endpoint Independent Mapping and Filtering (Full Cone NAT)
 - ALG's (fixups) should not be used (exceptions like A-FTP)
 - Paired IP address pooling behavior
 - Port Parity preservation for UDP
 - Hairpinning behavior
 - Static Port Forwarding (PCP)
- Management
 - Port Limit per subscriber
 - Mapping Refresh
 - Very scalable NAT logging (binary Netflow)
- Redundancy – Intra-box Active/Standby, Inter-box Active/Active
- Scale – 10M's concurrent sessions, 100K's sessions per second, Virtualization (VRF-aware)
- IPv6 Transition Tool-set – dual-stack, NAT64, 6RD, DS-Lite, MAP-T...

Paired IP Address Pooling Behavior



- Paired (recommended) : use the same external IP address mapping for all sessions associated with the same internal IP address
- Some peer to peer applications don't negotiate the IP address for multiple sessions (eg. apps that are not able to negotiate the IP address for RTP and RTCP separately)

Hairpinning Behavior



- Use Case: Allow communications between two endpoints behind the same NAT when they are trying each other's external IP addresses

Notation

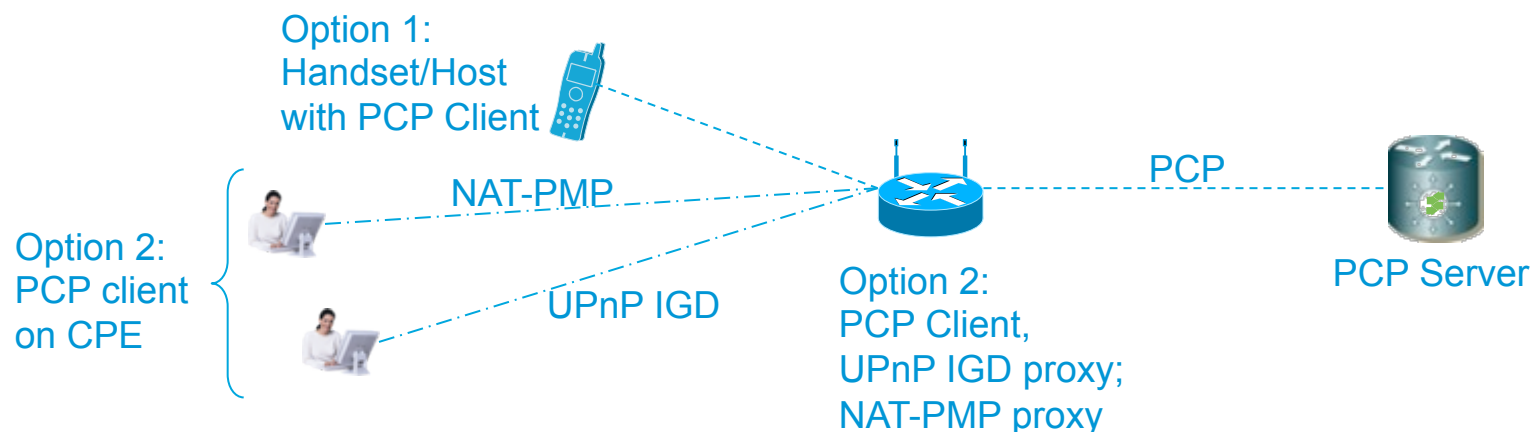
X:100

IPv4 address:Port *

* TCP/UDP port or Query ID for ICMP

Static Port Forwarding

- Requirement: Ability to configure, a fixed private (internal) IP address:port associated with a particular subscriber while CGN allocates a free public IP address:port
- New protocol: PCP (Port Control Protocol)



- Delegate port numbers to requesting applications/hosts to avoid requirement for ALGs
- [draft-ietf-pcp-base](#)

Other Port Behaviors

No Port Overloading

- A NAT must not have a "Port assignment" behavior of "Port overloading" (i.e. use port preservation even in the case of collision). Most applications will fail if this is used.

Port Parity Preservation

- An even port will be mapped to an even port, and an odd port will be mapped to an odd port. This behavior respects the [RFC3550] rule that RTP use even ports, and RTCP use odd ports.

Port Limit Per Subscriber

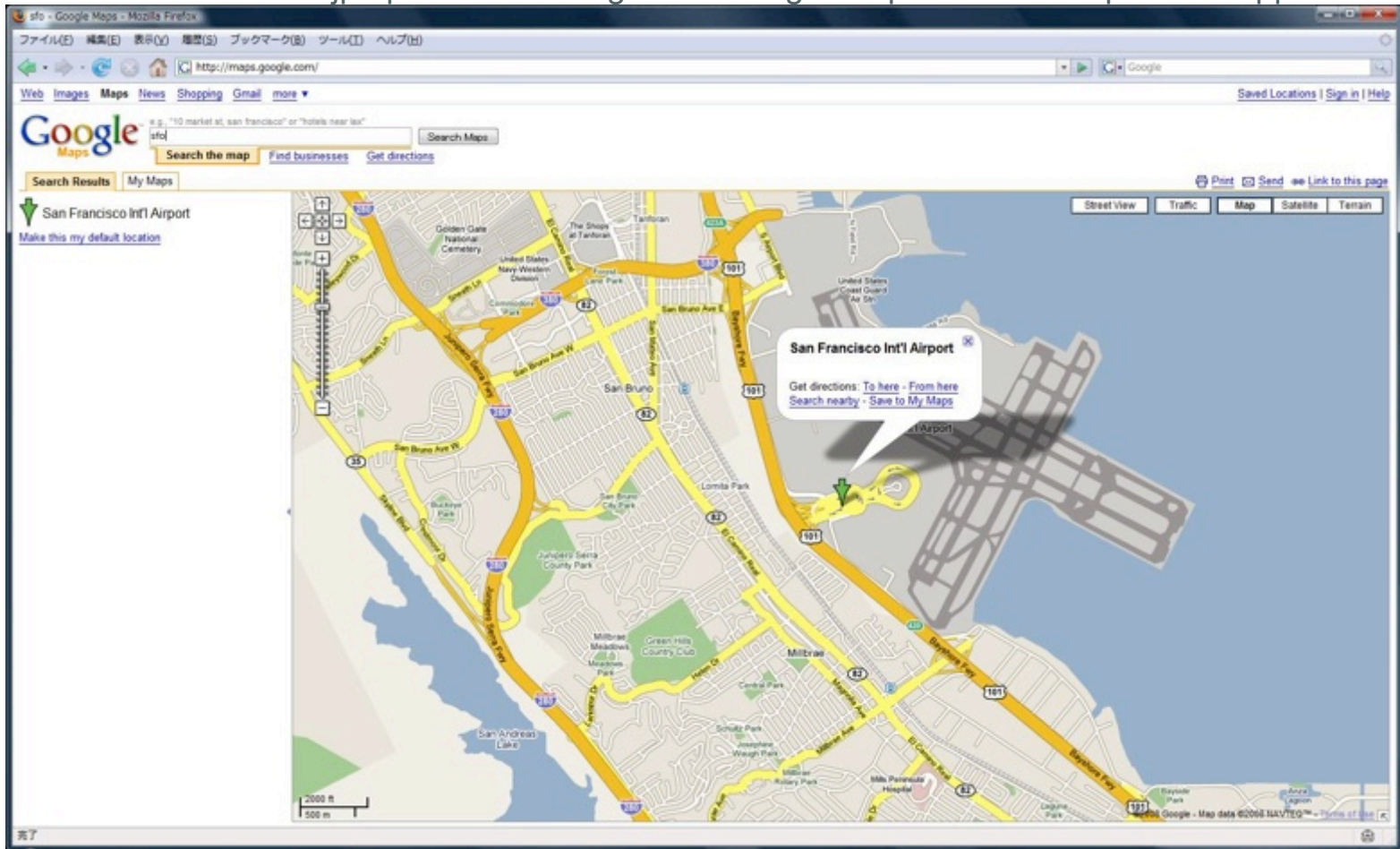
- Configurable port limit per subscriber for the system (includes TCP, UDP and ICMP). NAT Security – DoS attack/virus exhaust prevention.

Impact of NAT Port Limits

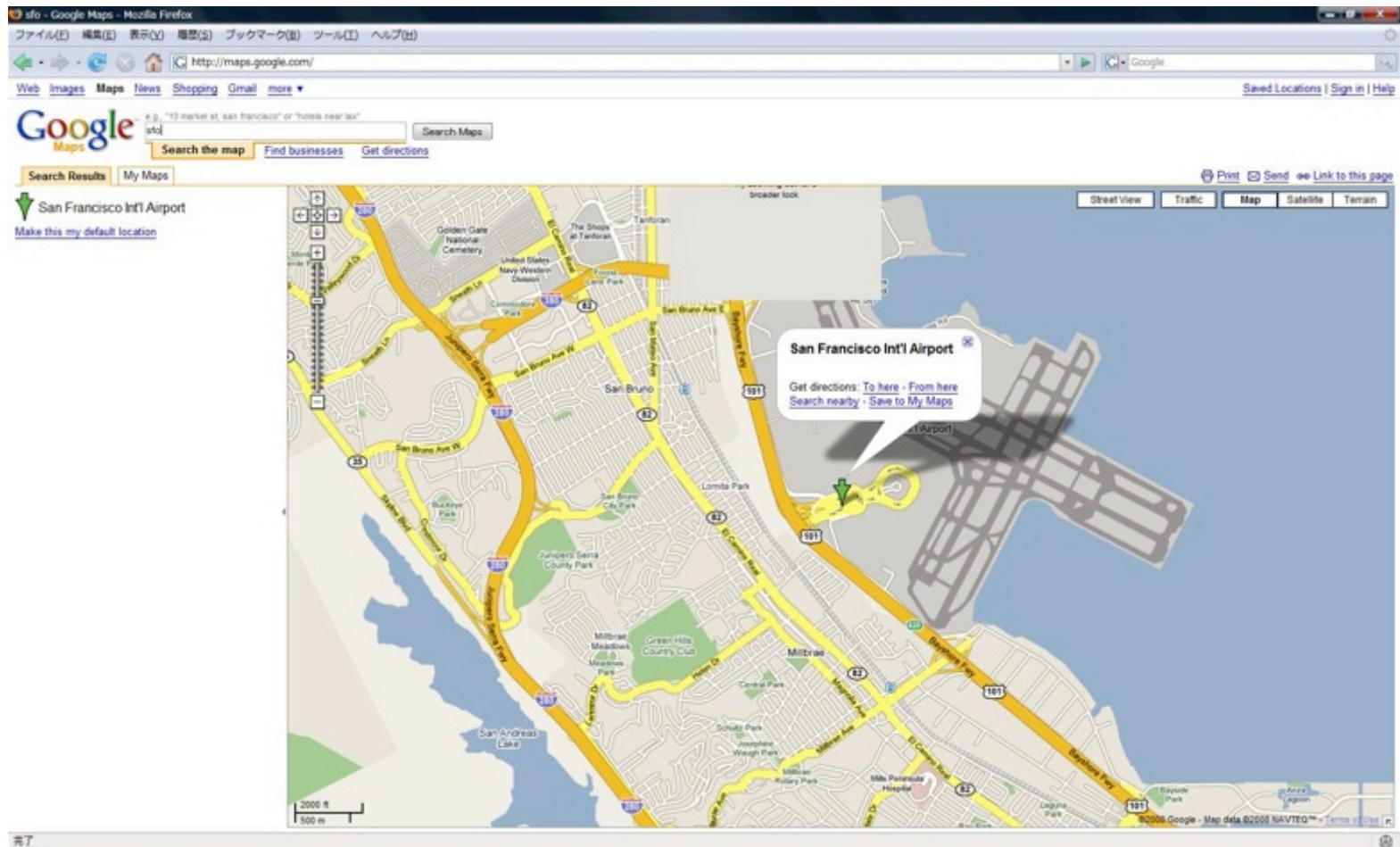
Example: GoogleMaps with Max 30 Connections

Example/Slides Courtesy of NTT, See Also:

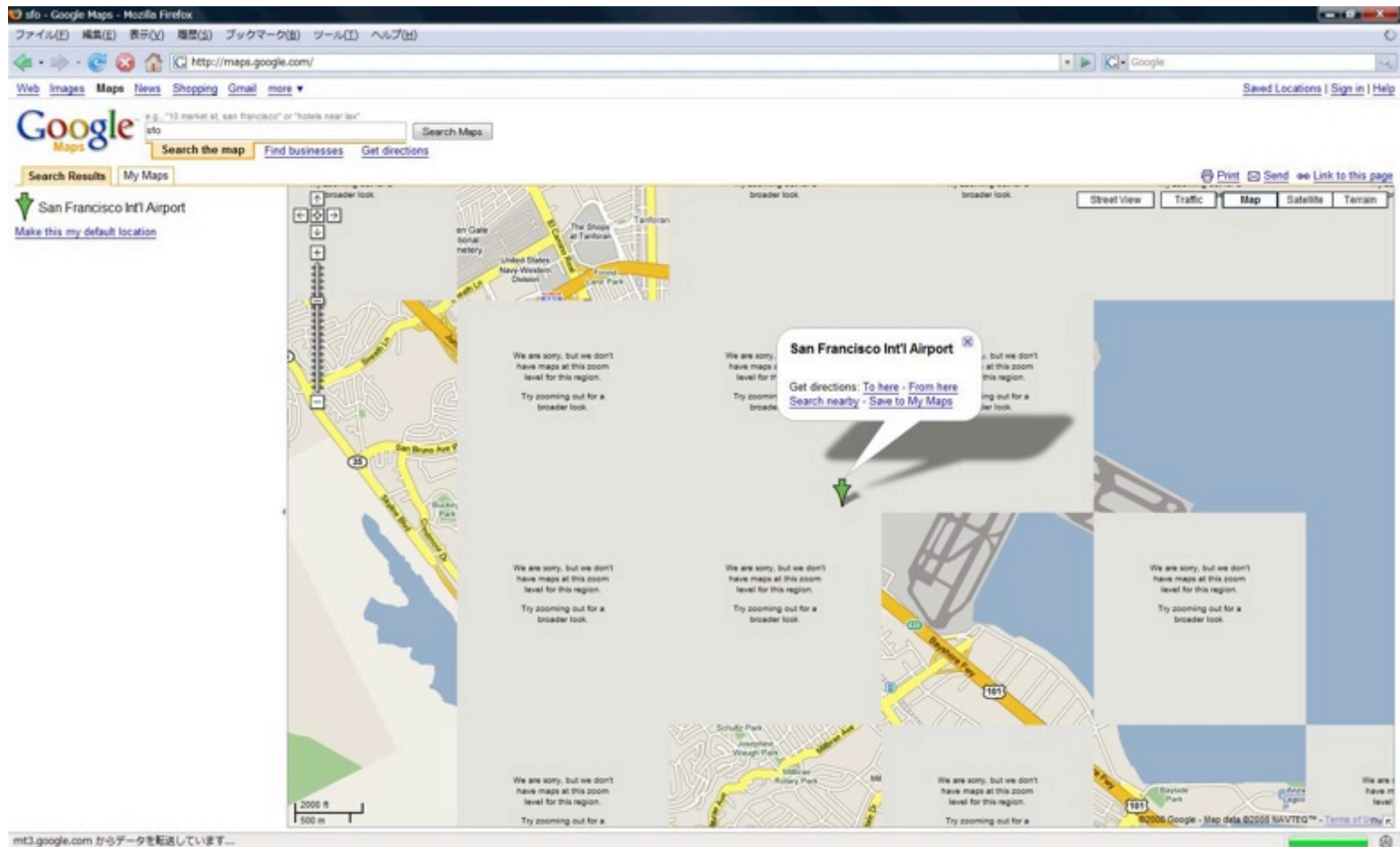
Hiroshi Esaki: www2.jp.apan.net/meetings/kaohsiung2009/presentations/ipv6/esaki.ppt



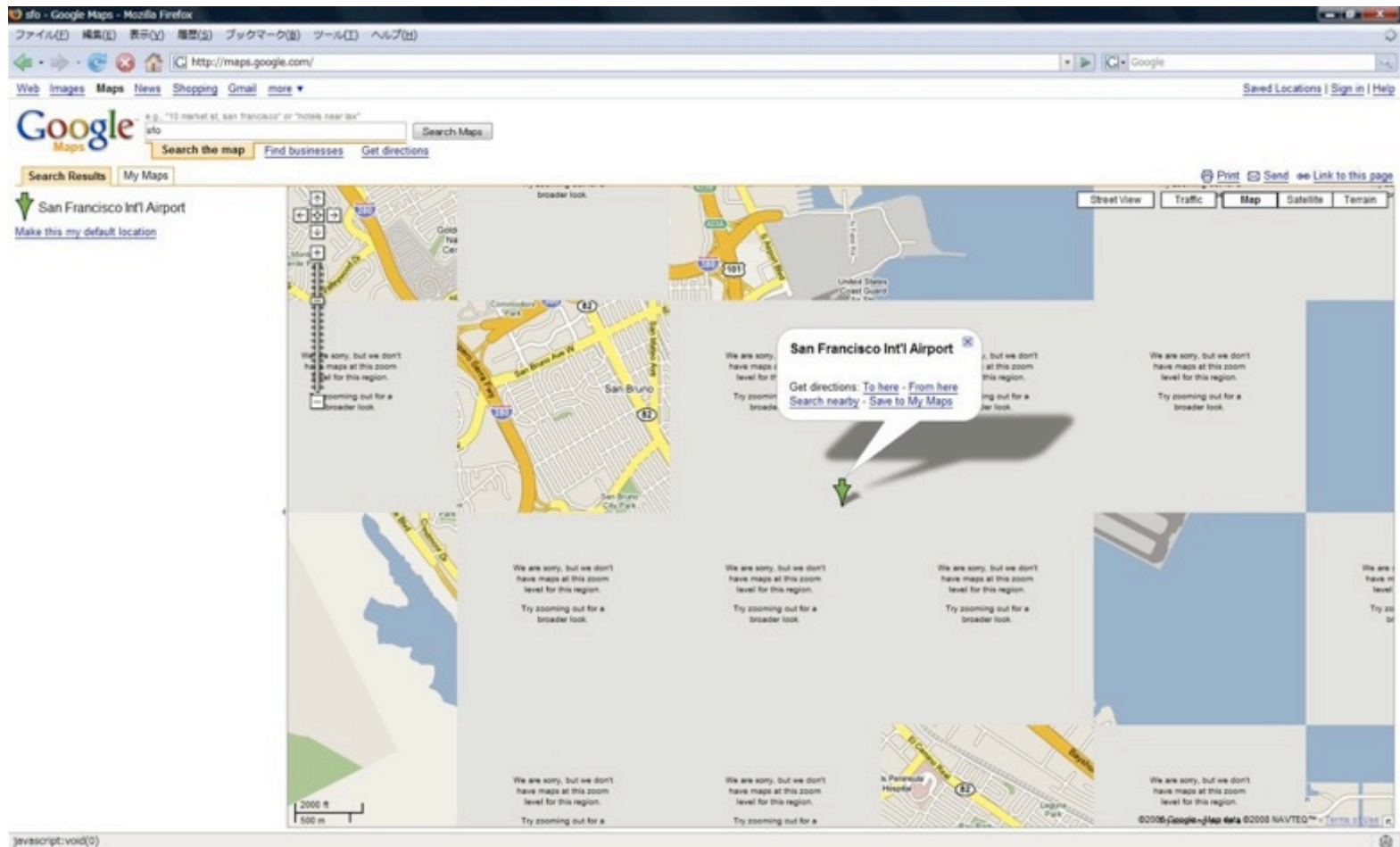
Max 20 Connections



Max 15 Connections



Max 10 Connections



Max 5 Connections



Number of Sessions for Some Applications

Some examples of major Web site

Application	# of TCP sessions
No operation	5~10
Yahoo top page	10~20
Google image search	30~60
ニコニコ動画	50~80
OCN photo friend	170~200+
iTunes	230~270
iGoogle	80~100
楽天(Rakuten)	50~60
Amazon	90
HMV	100
YouTube	90

* Courtesy of NTT, Hiroshi Esaki

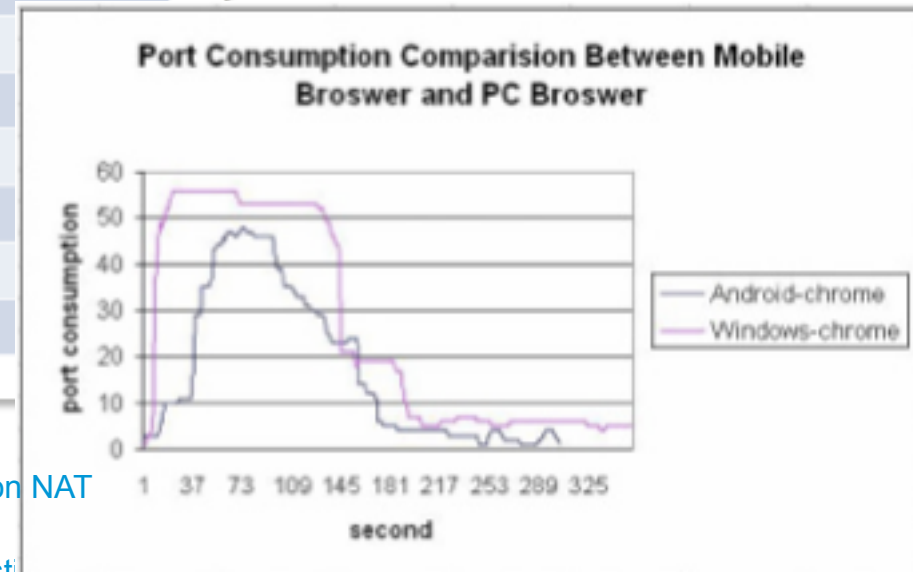
Source:

Application behaviors in in terms of port/session consumptions on NAT
<http://opensourceaplus.weebly.com/experiments-results.html>

See also “An Experimental Study of Home Gateway Characterist
<https://fit.nokia.com/lars/papers/2010-imc-hgw-study.pdf>
<http://www.ietf.org/proceedings/78/slides/behave-8.pdf>

Port Consumption can be big

- Eg. AJAX-based applications with tens-hundreds of TCP sessions
- Eg. Relaunching Firefox with Tabs opens hundreds of sessions



Default Session Timers – example

Type	Default Value
ICMP	60 sec
UDP init	30 sec
UDP active	120 sec
TCP Init	120 sec
TCP active	30 min

***) Default Refresh Direction is Bidirectional (configurable to OutBound only)**

CGN Session Logging

- Data Retention Law compliance, user trackability

Who posted a content to a server on Tue at 8:09:10pm?

- Global IP:port → CGN Log → Private IP:port → MSISDN
- [Directive 2006/24/EC - Data Retention](#)

- Logging Format

Must be fast and efficient (think 1Msps)

- ASCII format (Syslog) – very chatty (113B add-event), inefficient, no sequence #
- Binary Format (Netflow) – efficient (21B add event), sequencing

- Netflow V9 Logging

21B add-event, 11B delete-event

Up to 68 add-events per 1500B export packet

- Dynamic, template-based format (1 Msps = cca 176 Mbps, 14.7 Kpps)

Future evolution → IPFIX (SCTP – reliable streaming, multi-core CPU's)

Netflow v9 logging

Tip: IsarFlow – tested CGN NFv9 Collector
www.isarnet.de

Add Event
Template 256
(21B)

Field ID	Attribute	Value
234	Incoming VRF ID	32 bit ID
235	Outgoing VRF ID	32 bit ID
8	Source IP Address	IPv4 Address
225	Translated Source IP Address	IPv4 Address
7	Source Port	16 bit port
227	Translated Source Port	16 bit port
4	Protocol	8bit value

Delete Event
Template 257
(11B)

Field ID	Attribute	Value
234	Incoming VRF ID	32 bit ID
8	Source IP Address	IPv4 Address
7	Source Port	16 bit port
4	Protocol	8bit value

Netflow logging Data Volume example

Collector Performance – 100K users, average and peak

Average NF Collector Load	
NF records per sec	188.800
Avg Load - Logging BW Rate (Mbps)	33
Avg Load - Logging Packets per Second (pps)	2.776
Peak NF Collector Load during CGSE Switchover	
Peak Load - NF records per sec	1.000.000
Peak Load - Logging BW Rate (Mbps)	176
Peak Load - Logging Packets per Second (pps)	14.706
Peak Load - Duration (s)	9

Storage Capacity – includes per-day user behavior

Storage Data Volume	
Average Subscriber active Duration (hrs/day)	8
Raw NF Data Volume per Day (Gbytes)	29
Total Data Volume per year (TBytes)	10.5
Database Overhead Factor	2.0
Database Compression Rate	6.0
Total Data Volume per year (TBytes) - compressed	3.5

Reality check: 100K CGN users would consume 3.5TB storage per year
(compressed, fully SQL searchable data)

E-Shop: 4TB disk, 300 Euro...

Usually no need to bother with logging reduction...

Logging reduction

- Bulk port range allocation
 - Pre-allocates a port-set per user (eg. 512 ports)
 - Logs only once (when port-set is allocated/deallocated)
- Deterministic NAT
 - Port-sets are determined algorithmically from user IP (patented)
 - No logging until port-set overflows

ISSUES

- Breaks TCP port randomization (user security consequences)
- Inefficient usage of global IP pools
(active users eat hundreds of ports, inactive users eat only few ports)
- Troubles to get acceptance in IETF (BEHAVE WG)
- Mutually exclusive with DBL (Destination Based Logging)

→ It is not worth it and is **NOT RECOMMENDED**. It's better to use Netflow.

DBL (Destination Based Logging)

Add Event
Template 271
(27B)

Field ID	Attribute	Value
234	Incoming VRF ID	32 bit ID
235	Outgoing VRF ID	32 bit ID
8	Source IP Address	IPv4 Address
225	Translated Source IP Address	IPv4 Address
7	Source Port	16 bit port
227	Translated Source Port	16 bit port
12	Destination Address	IPv4 Address
11	Destination Port	16 bit port
4	Protocol	8 bit value

DBL logs also destination IP:port

- data retention vs. user privacy
- keeps EIM/EIF behavior

CGN Scale and Performance

Session = full-duplex, bidirectional L4 flow

- Session Setup Rate [sps] – sessions per second

Average # of New Sessions per User, during peak hours

- Huge load during a failover scenarios or after a power blackout
- Failing to cope with SPS = huge TCP delays, timeouts/retransmissions

- Maximum Number of **Concurrent Sessions [cs]** per CGN

Average # of Concurrent Sessions per User, during peak hours

- UDP must not expire in less than 2 minutes (RFC4787)
- UDP/TCP timers for Initializing and Established sessions should be configurable

- **Throughput** per CGN [bps]

Aggregate (downstream + upstream) bandwidth

CGN Scale and Performance – design

L (Low-scale) Scenario – 3G mobile users, smart-phones

M (Medium-scale) Scenario – ADSL subscribers, PC users with 3G/4G dongles, Tablets, WiFi and top smart-phone users

H (High-scale) Scenario – heavy Broadband users, Internet sharing

	H	M	L
Average BW – Σ of combined DL + UL BW per subscriber during peak <u>hrs</u>	300 kbps	150 kbps	30 kbps
Average #of Concurrent Sessions – Σ of all ports/protocol and direction	200	100	10
Maximum #of Concurrent Sessions – Σ of all ports/protocol and direction	1000	500	100
Average Session Transaction Rate – #of inbound and outbound sessions that are established or deleted per second	4,0	1,0	0,1

100K BB users = up to 100Ksps and 10Mcs during peak hour!

CGN Redundancy

- High Availability scenarios

Intra-chassis, Inter-chassis

Active/Standby, Active/Active

- Stateful or stateless

Millions of short-lived Layer-4 session

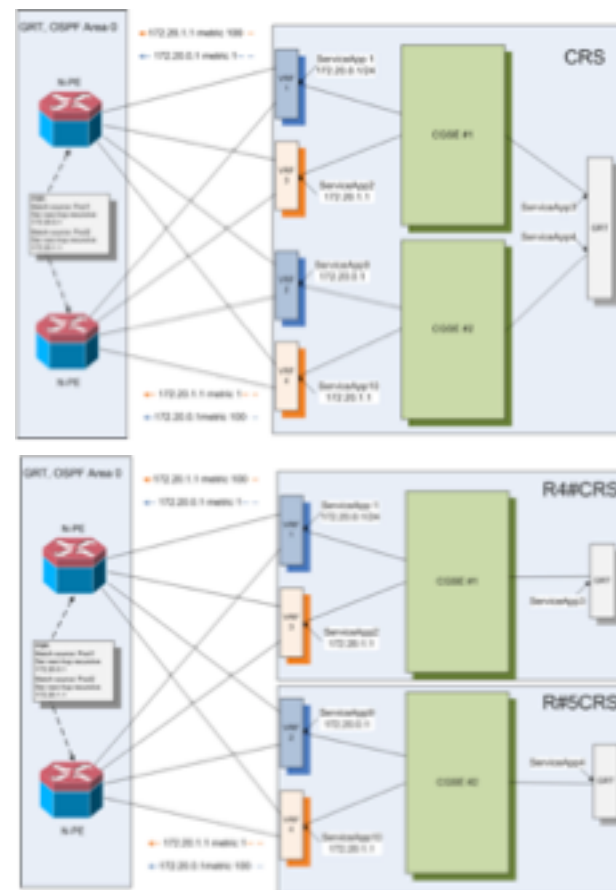
Stateful sync makes no sense for such ephemeral state (memory & CPU) – eg.

- Stateless redundancy

1Msps = 100K active users (10Mcs) are up in 10s → minimal loss

Load-sharing = simple ECMP routing

Best Practice: Simple Non-Revertive 1:1 Warm Standby



CGN design: Basic Scenarios

CGN module

1+1 Redundancy

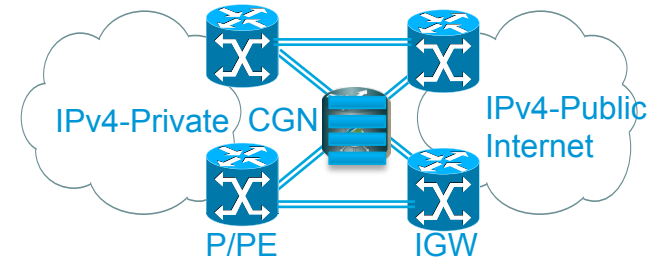
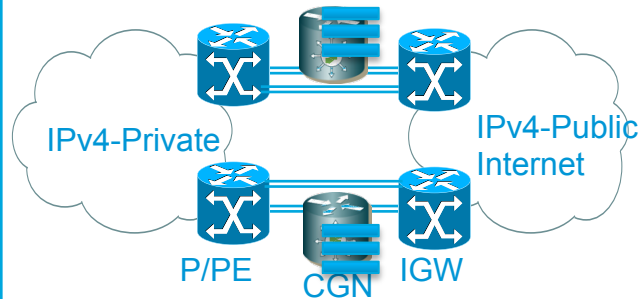
- 3+3 in this example

N+1 Redundancy

- 3+1 in this example

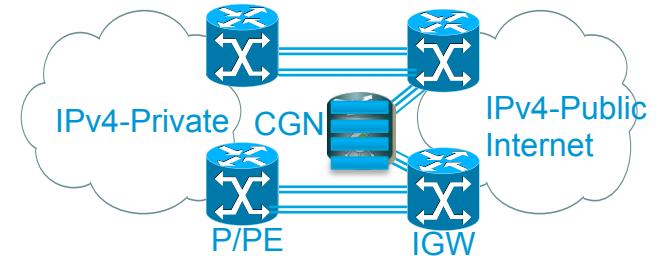
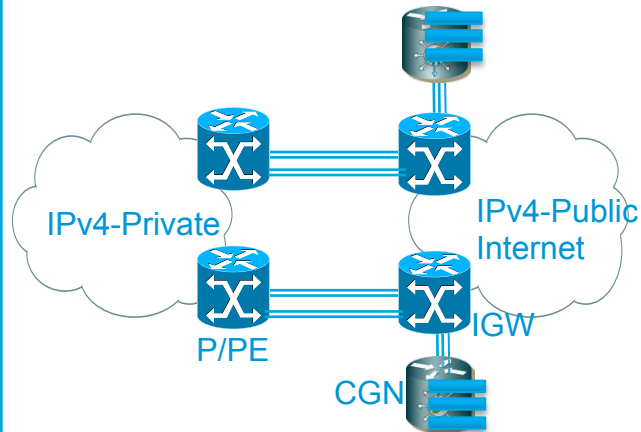
Bump-in-a-wire Design

- new 4xTGE per box in this example



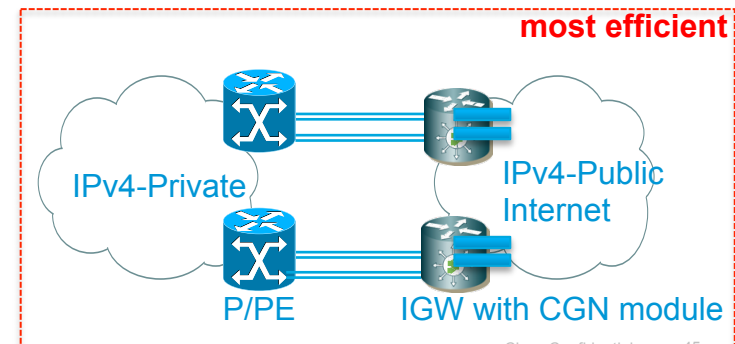
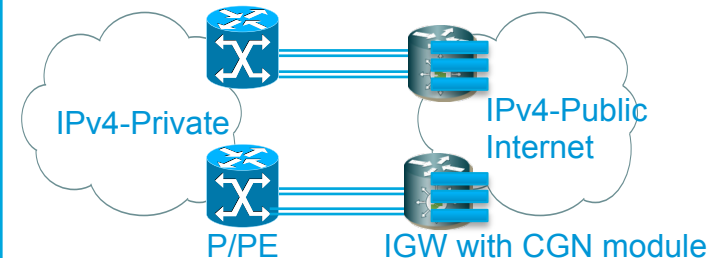
Router-on-a-stick Design

- new 3xTGE per box in this example



Integrated Design

- No new ports



Summary

- CGN is here to overcome IPv4 exhaust before IPv6 migration
- CGN is ISP element, focus on transparency
- CGN is not firewall
- CGN behavioral requirements
- CGN logging
- CGN performance – SPS, # of sessions, deployment options

Thank you.

